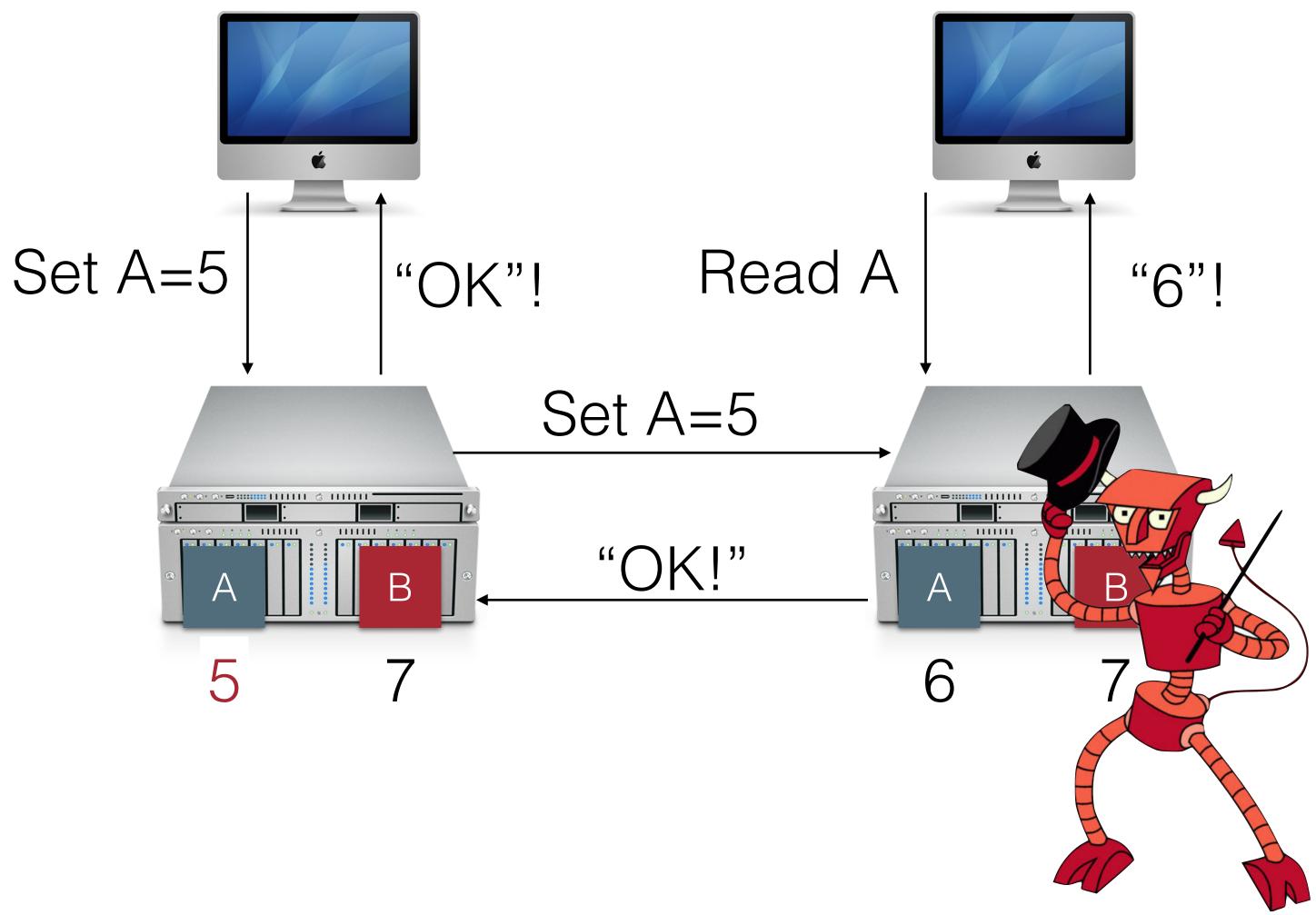
### Security in Distributed Systems

CS 475, Spring 2019 Concurrent & Distributed Systems



### **Byzantine Faults**





#### Blockchains

- Solution: make it hard for participants to take over the network; provide rewards for participants so they will still participate
- Each participant stores the entire record of transactions as blocks
- Each block contains some number of transactions and the hash of the previous block
- All participants follow a set of rules to determine if a new block is valid

$$h_0 \bigcup_{i=1}^{n_0} h_1 \bigcup_{i=1}^{n_0} h_2 \bigcup_{i=1}^{n_0} h_3 \bigcup_{i=1}^{n_1} h_4 \bigcup_{i=1}^{n_1} h_6 \bigcup_{i=1}^{n_0} h_7 \bigcup_{i=1}^{n_1} h_8 \bigcup_{i=1}^{n_0} h_n \begin{bmatrix} d_n \end{bmatrix}$$



- Today:
  - Security in Distributed Systems
  - Discussion of course structure
- Final Exam info:
  - $\bullet$ granularity)
- Reminder Project is out!
  - Fault-tolerant, sequentially consistent replicated key value store
  - Can do in a group (1 to 3 students per group)



#### Primary focus is second half of semester (starting from Networks), still responsible for concepts from first half (notably mutual exclusion, locking



## Security isn't (always) free

- do to protect your belongings/property?
- Do you change the locks?
- Do you buy security cameras?
- Do you hire a security guard?
- Do you even bother locking the door?

• You just moved to a new house, someone just moved out of it. What do you



## Security: Managing Risk

- Security architecture is a set of mechanisms and policies that we build into our system to mitigate risks from threats
- Threat: potential event that could compromise a security requirement
- Attack: realization of a threat  $\bullet$
- Vulnerability: a characteristic or flaw in system design or implementation, or in the security procedures, that, if exploited, could result in a security compromise





#### What does it mean for a distributed system to be secure?

- Maintain a secure channel between nodes:
  - Authenticity (Who am I talking to?)
  - Confidentiality (Is my data hidden?)
  - Integrity (Has my data been modified?)
  - Availability (Can I reach the destination?)
- Maintain some security about who participates in the system?
- What cryptographic tools are available to us?





### Costs & Benefits

- Increasing security might:
  - Increase development & maintenance cost
  - Increase infrastructure requirements
- Degrade performance
- But, if we are attacked, increasing security might also:
  - Decrease financial and intangible losses
- So: How likely do we think we are to be attacked in way X?

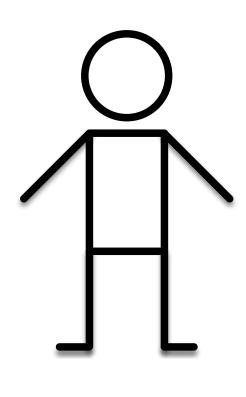


### Threat Models

- What is being defended?
  - What resources are important to defend?
  - What malicious actors exist and what attacks might they employ?
- Who do we trust?
  - What entities or parts of system can be considered secure and trusted
  - Have to trust **something**!

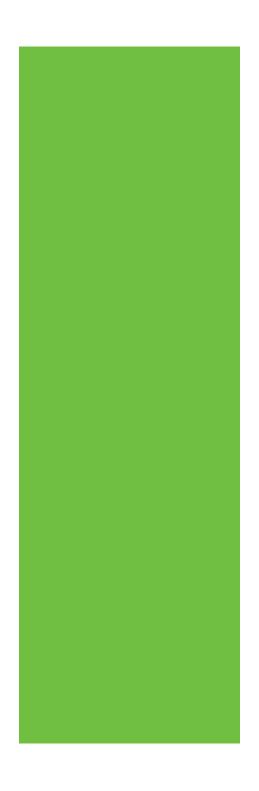


**HTTP Request** 



**HTTP Response** 

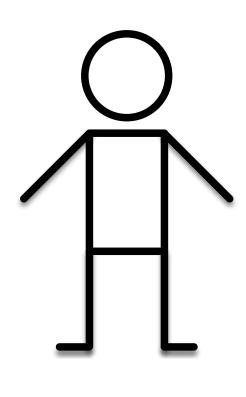
#### client page (the "user")



#### server

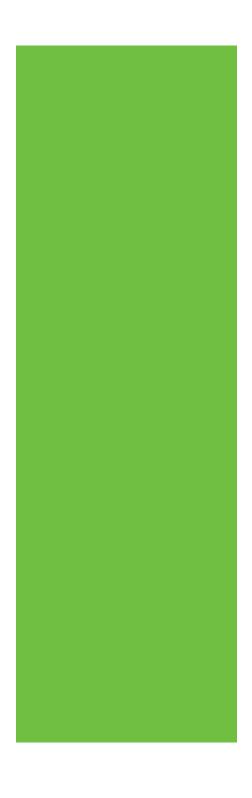


**HTTP Request** 



**HTTP Response** 

#### client page (the "user")



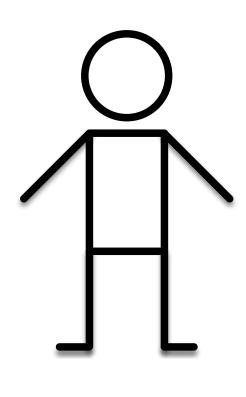
#### server

#### **Do I trust that this request** *really* came from the user?

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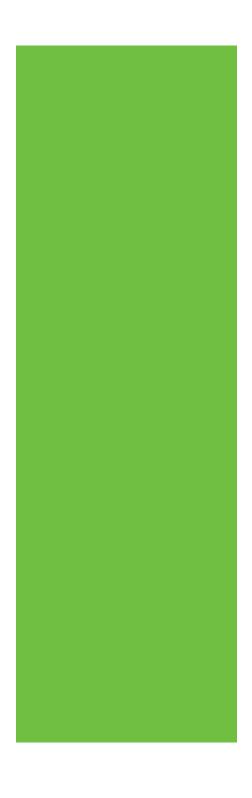
**HTTP Request** 



**HTTP Response** 

#### client page (the "user")

**Do I trust that this response** *really* came from the server?



#### server

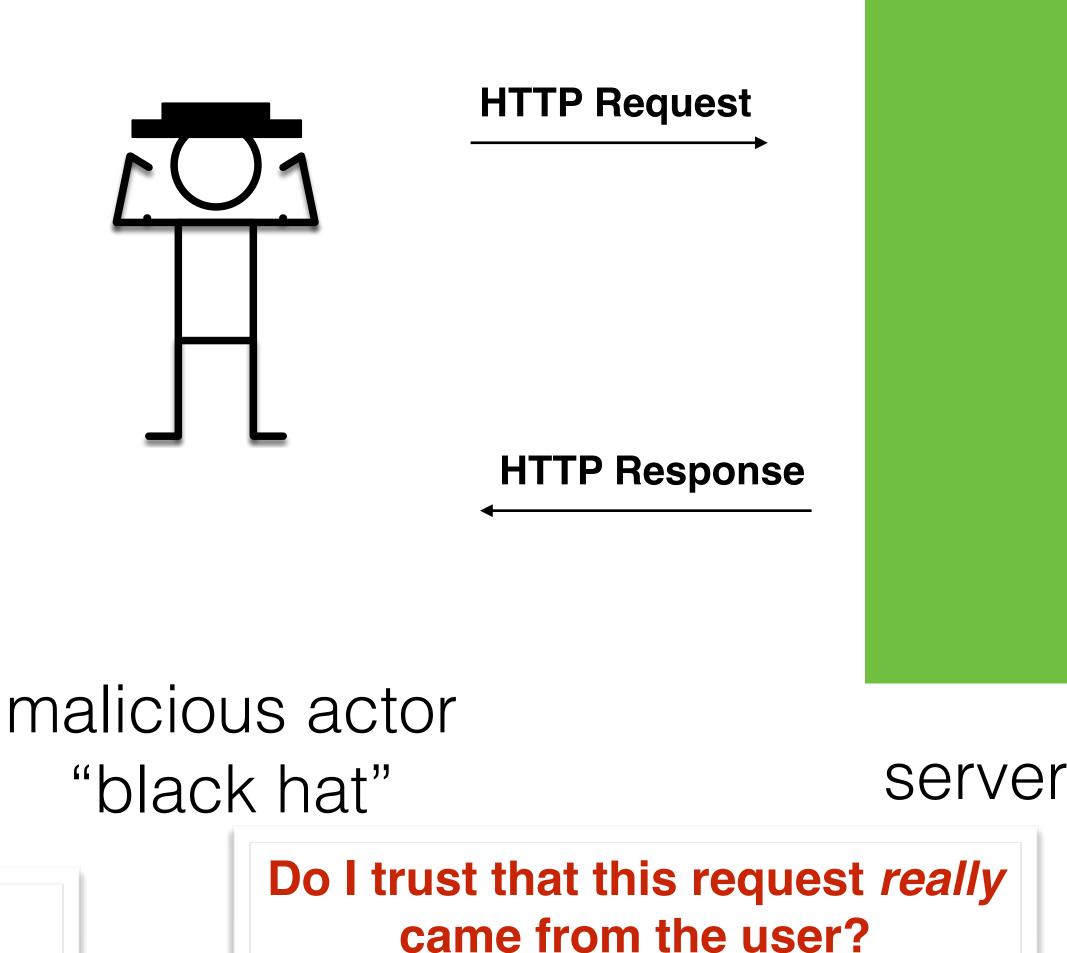
#### **Do I trust that this request** *really* came from the user?



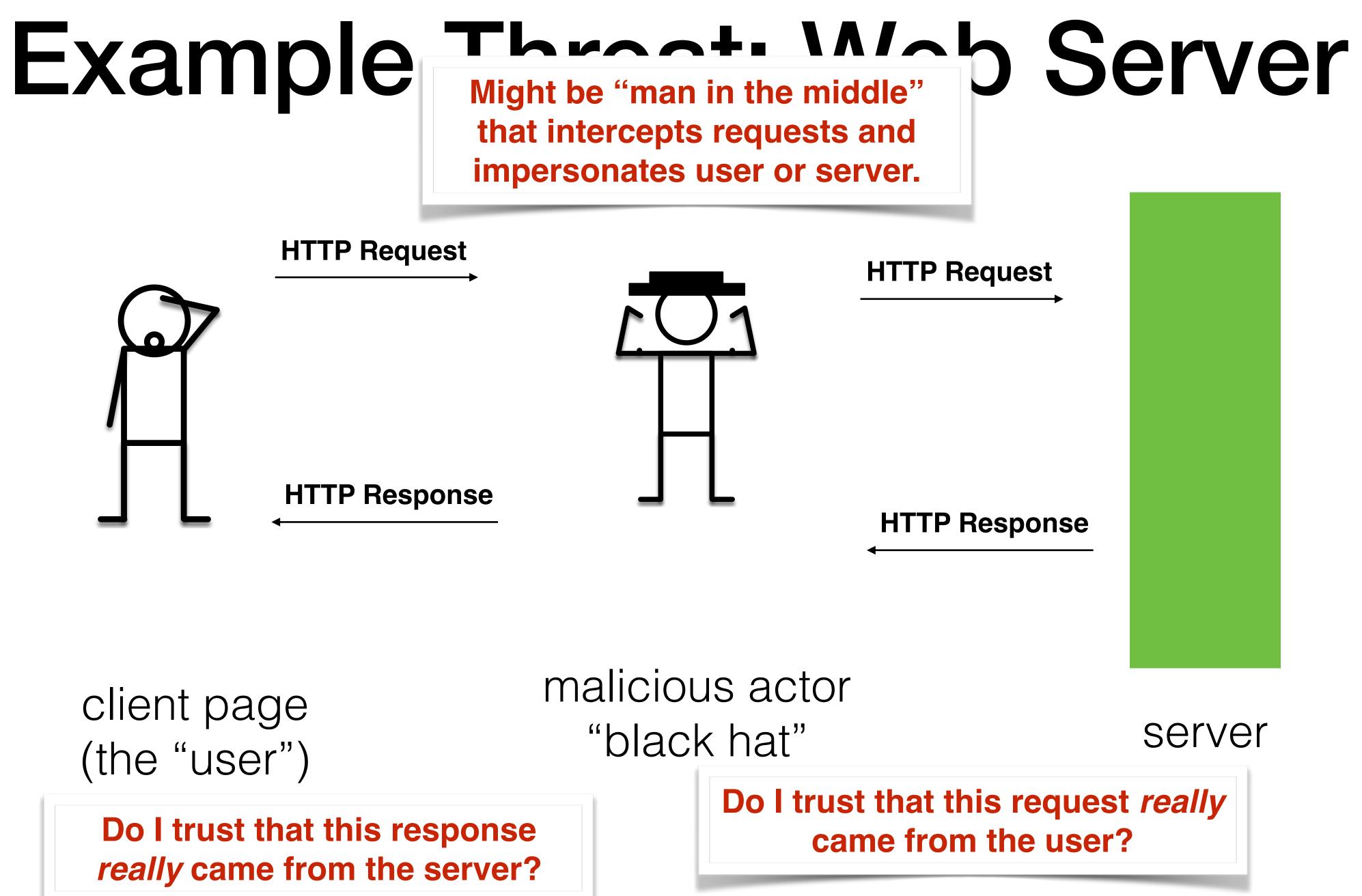
**HTTP Request HTTP Response** 

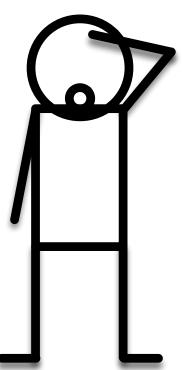
client page (the "user")

**Do I trust that this response** *really* came from the server?











#### **Other Risks**

- Is our network well behaved?
- Is our network malicious?
- Who can access our system?
- Are our users well-behaved?
- Are our users malicious?
- Is our system well behaved?



### **Protection Concerns**

- Secure channels of communication
  - Authentication: is everyone who they say they are?
  - Confidentiality & integrity: is a third party interfering in our communication?
- Access Controls
  - Authorization: Who has access to an operation/resource?
  - Accountability: Maintaining an audit trail
  - Non-repudiation: A participant can not deny some action that they took with the system

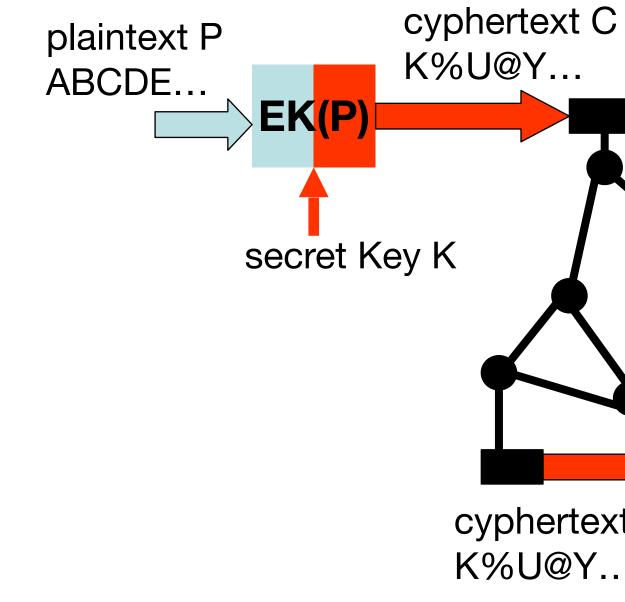






#### Symmetric encryption, aka shared secret key

- $\bullet M = D_{K}(E_{K}(M))$
- •M is the data, D is decrypt, E is encrypt, and k is the key
- •Computationally efficient (relatively)
- Can have hardware support too (e.g. iPhone)



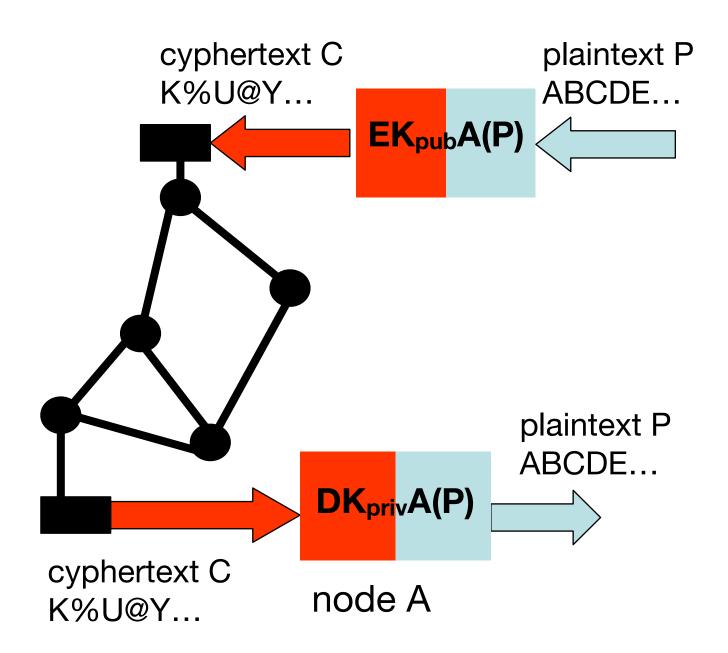
DK(C plaintext P ABCDE... cyphertext C K%U@Y... secret Key K





#### Asymmetric encryption, aka public key/private key

- $M = DK_{priv} (EK_{pub} (M)) = DK_{pub} (EK_{priv} (M))$
- When a node B wants to send a message to node A, it obtains A's public key and uses K<sub>APublic</sub> to encrypt the message
- Only A can decrypt the message using its private key
- Computationally expensive









## **Public/Private Key Encryption**

• Encrypt with public key: only private key holder can decrypt

Public Key

Plain text Message

Private Key

Encrypted Message

Plain text Message



## **Public/Private Key Encryption**

• Encrypt with private key: anyone with public key can decrypt

Private Key

Plain text Message

Public Key

Signed Message

Plain text Message



## Hashing

- S = H(M)
- intentional change to the data will change the representation
- Fixed size and independent of size of M
- Computationally efficient

#### Inputs

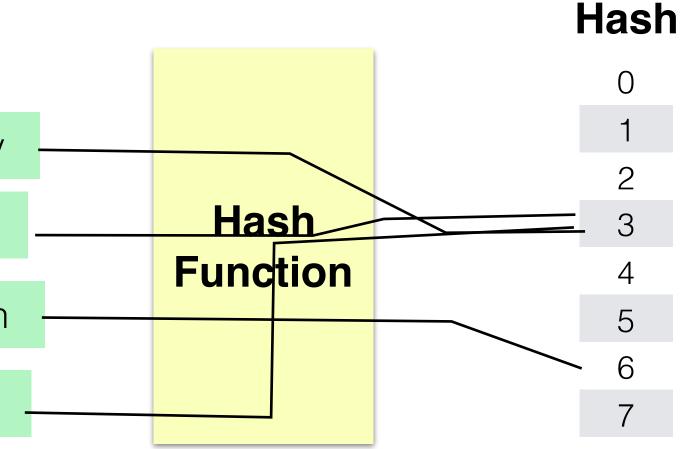
Leo McGarry

Josh Lyman

Sam Seaborn

Toby Ziegler

# S, aka digest, is a unique representation of data such that an accidental or





## Hashing to verify messages

- Just sending a hash of the message isn't enough!
- How do we know that a third party didn't tamper with the message and the hash?
- Solution: encrypt the hash using your private key. Anyone can verify that the hash was "signed" by you

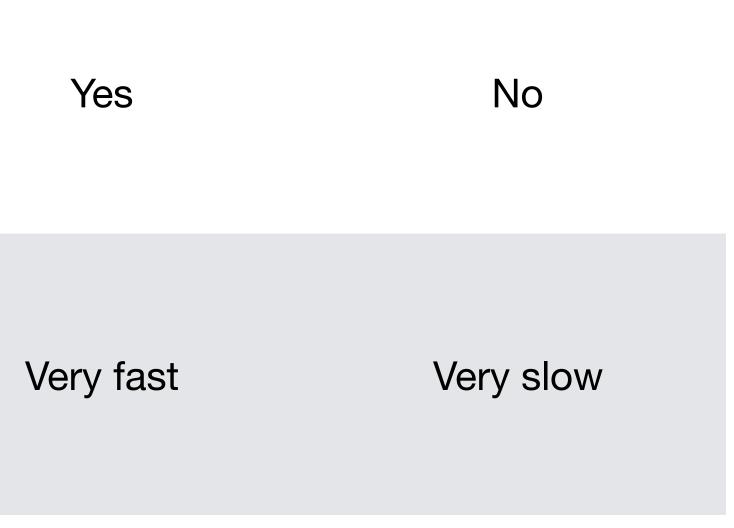


### Symmetric vs Asymmetric Crypto

**Requires a pre**shared secret

**Relative speed** 

Symmetric Crypto Asymmetric Crypto







## Asymmetric Cryptography

- So, great: no need to pre-share anything, right!
- Widely used for instance... HTTPS! SSL!
- But: there's a bootstrapping problem  $\bullet$

- When you visit <u>amazon.com</u>, the site will sign its content using its private key • You can use <u>amazon.com</u>'s public key to verify it's really from amazon.com How do you know what <u>amazon.com</u>'s public key is though?
- "PKI" Public Key Infrastructure

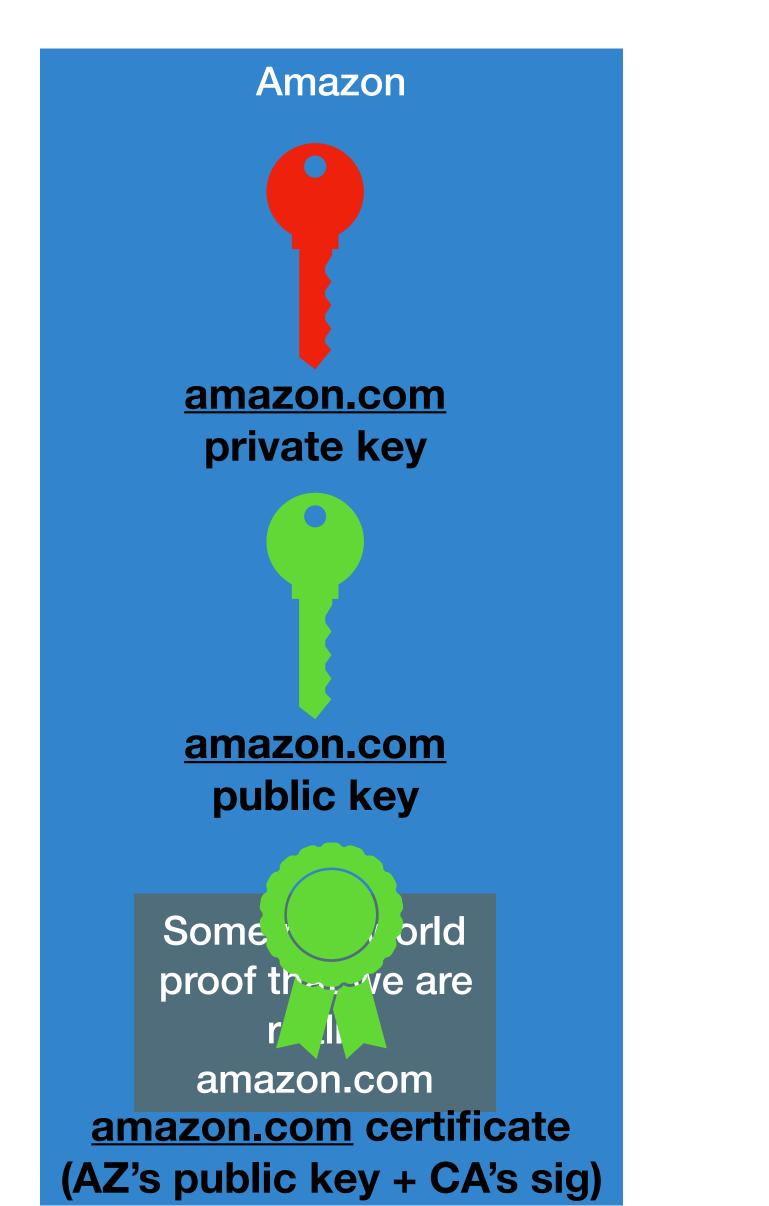




- A certificate authority (or CA) binds some public key to a real-world entity that we might be familiar with
- The CA is the clearinghouse that verifies that <u>amazon.com</u> is truly amazon.com
- CA creates a certificate that binds <u>amazon.com</u>'s public key to the CA's public key (signing it using the CA's private key)







J. Bell

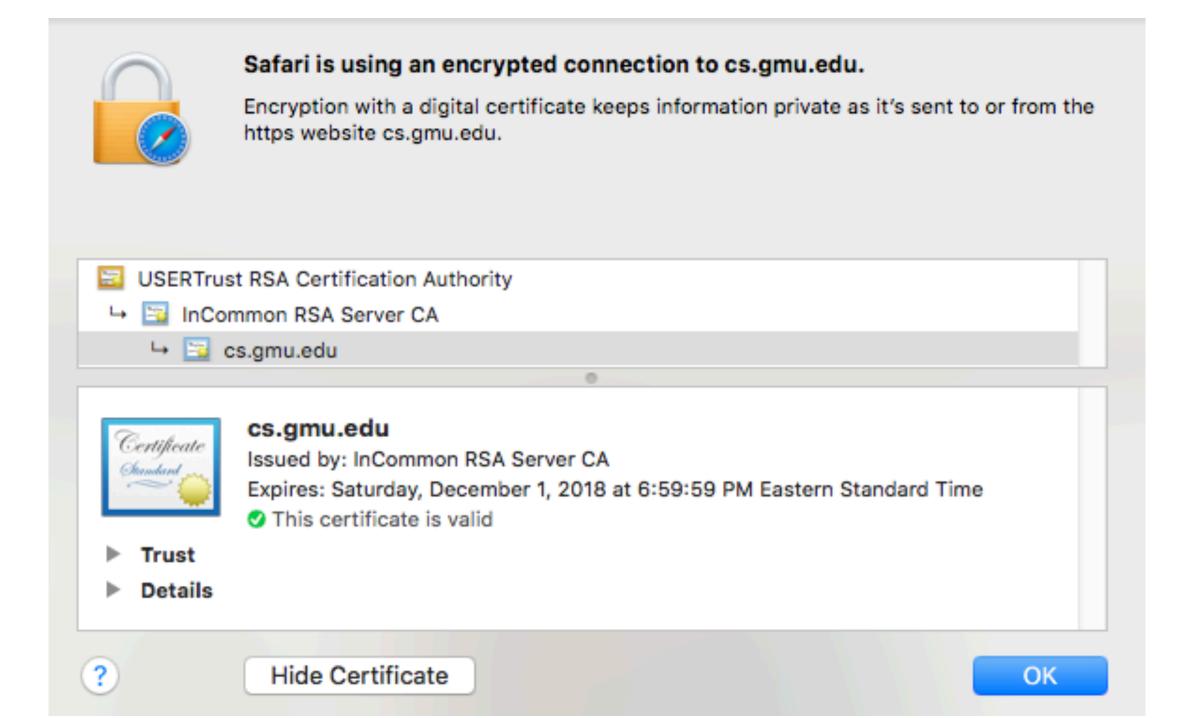
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amazon.com certificate (AZ's public key + CA's sig) **Certificate Authority** CA private key **CA public key** 





- Note: We had to already know the CA's public key
- There are a small set of "root" CA's (think: root DNS servers)
- Every computer/browser is shipped with these root CA public keys



#### CA's public key (think: root DNS servers) d with these root CA public keys



- What happens if a CA is compromised, and issues invalid certificates?
- Not good times.

#### Security

#### **Comodo-gate hacker brags about** forged certificate exploit

Tiger-blooded Persian cracker boasts of mighty exploits

#### Security

#### Fuming Google tears Symantec a new one over rogue SSL certs

We've got just the thing for you, Symantec ...

By Jain Thomson in San Francisco 29 Oct 2015 at 21:32 36 🖵 SHARE V



Google has read the riot act to Symantec. scolding the security biz for its



### **Denial of Service Attacks**

- A significant concern for distributed systems
- An attack on **availability** attackers prevent legitimate users from accessing system
- Can attack:
  - Bandwidth
  - CPU
  - Memory
- Core problem:
  - Costs more to process a message than to send it





#### **Distributed Denial of Service Attacks (DDoS)**

- Model: Attacker has (hundreds of?) thousands of machines at disposal to attack
- Most common form of DoS today
- Exhausts network bandwidth
- Typically rooted in a botnet some command and control infrastructure setup by an attacker, who then controls all of these machines







#### Strawman Defenses

- Make a filter list of bad addresses?
- Trace down the person responsible?



### Heuristic Defenses

- Overprovision
- Black-hole routing
- Filter anomalies
- Replication



### Overprovisioning

- Make a DDoS-proof site by making it far bigger than it needs to be
- expect
- A losing battle: an attacker can always get more bots!

• Provision 100x bandwidth, 100x server capacity etc. compared to what you



### Black-Hole Routing

- Limits the impact of an attack
- ISP re-routes traffic to the target site to a black hole
- Site still goes offline
- But not crashed, other sites on servers sharing network links are OK
- Most DDoS attacks are short-lived, so clears up later



### Anomaly Filtering

- DDoS traffic usually has something peculiar about it...
  - Automatically generated requests following a pattern? lacksquare
- Route all traffic through black-box filters that try to learn this stuff and identify anomalies
- Imperfect, but often works





#### Other DoS attacks

- Reflector
- Complexity



### **Reflector Attacks**

- Exploits a publicly available service to amplify an attack
- Example: DNS
- Attacker makes a (relatively small) DNS request  $\bullet$
- Attacker forges their own IP address with the victim's
- DNS server responds to the victim's IP address



### **Complexity Attacks**

- Increasingly common as we find defenses for other attacks
- Idea: Can I make one request that is 100 times as hard to process as other requests?
  - Then I only need to make 1% of the requests I would have had to otherwise, in order to get the same attack!





#### **Billion** Iolz

```
<?xml version="1.0"?>
<!DOCTYPE lolz [
<!ENTITY lol "lol">
<! ELEMENT lolz (#PCDATA)>
]>
<lolz>&lolz;</lolz>
```

#### After parsing: this document contains "lol" repeated literally a billion times... ~3GB of RAM

<!ENTITY lol2 "&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;&lol1;\*</pre> <!ENTITY lol3 "&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;&lol2;\*> <!ENTITY lol4 "&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;&lol3;</pre> <!ENTITY lol5 "&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;&lol4;"> <!ENTITY lol6 "&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;&lol5;</pre> <!ENTITY lol7 "&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;&lol6;</pre> <! ENTITY 1018 "&lol7; &lol7; "> <!ENTITY lol9 "&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;&lol8;</pre>



### **Discussion of course structure**

- First assignment without concurrency?
- More/less concurrency discussion at start?
- More/less programming?
- More/less detail on some topics?



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